Phase Transitions in Quantum Games

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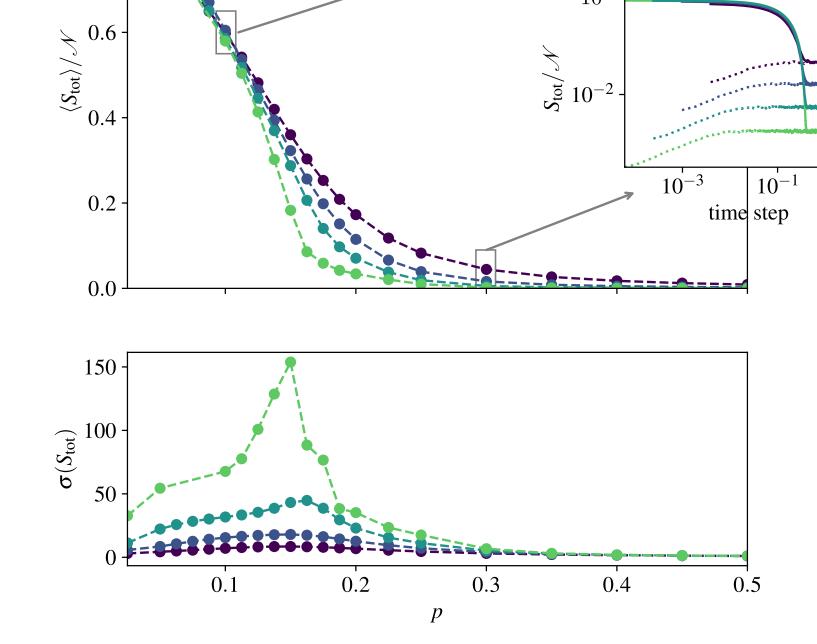
For each bond $i \in [1, N-1]$ Game set-up Motivation we compute the entanglement entropy • Random quantum circuits display interesting and universal behaviours, $S(i) = -\operatorname{Tr}[\rho' \log \rho']$ Two agents playing against each other 1 providing a simple tool to probe complex quantum dynamics • Game-like settings are highly tunable setups that reveal new and intriguing Entangler physical phenomena plays with probability 1-p1) selects randomly a bond $i \in [1, N-1]$, • Entanglement transitions, from area-law to volume-law steady states, have 2) samples a Clifford gate $g \in \mathcal{C}_2$, been observed in both mesurement-doped circuits⁶, but most recently also 3) acts with g on ith bond. in unitary gates circuits¹ • The degree and nature of accessible information significantly impact the physics, allowing us to modify entanglement transitions Disentangler plays with probability p 1) selects a bond $i \in [1, N-1]$ with policy, 2) chooses optimally disentangling Clifford gate g^* , 3) acts with g^* on ith bond. $policy \in \{ \text{ random, greedy, RL } \}$ At each time step Disentangler or Entangler N qubits plays according to p

Clifford circuits in a Clifford group \mathcal{C}_n : the group of unitaries that normalize the Pauli group (n = # qubits) $\mathcal{C}_n := \{ c \in U(2^n) \mid c\mathcal{P}c^{\dagger} = \mathcal{P}_n \}$ Pauli group $\mathcal{P}_n := \{\alpha O_1 O_2 ... O_N\},\$ with $\alpha \in \{1, -1, i, -i\}, O_j \in \{I, X, Y, Z\}$ \hookrightarrow The elements of \mathcal{C}_n are called **Clifford gates**, and are generated by *Hadamard* (H), *Phase* (S) and CNOT gates. Why Clifford circuits?^{2,3} $S = \begin{pmatrix} 1 & 0 \\ 0 & i \end{pmatrix}$ The Gottesman-Knill theorem ensures that any Clifford circuit can be simulated on a classical machine in time poly(n). CNOT = \rightarrow Can simulate larger system sizes and get closer to the thermodynamic limit **idea**: describe state $|\psi\rangle$ with its *stabilizers* instead of the amplitudes. \hookrightarrow unitary operator s.t. $s_i |\psi\rangle = |\psi\rangle$ \rightarrow stabilizer group Stab $(|\psi\rangle)$ Example State: $|\psi\rangle = |11\rangle$ stabilized by II, ZI, IZ, ZZ $x_{21} \quad \dots \quad x_{2N} \mid z_{21} \quad \dots \quad z_{2N}$ Tableau $\mathcal{T} =$ $1 \left(\begin{array}{cc|c} 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{array} \right)$ \Rightarrow tableau ($\left| \left\langle x_{N1} \quad \dots \quad x_{NN} \mid z_{N1} \quad \dots \quad z_{NN} \right\rangle \right|$ since $\operatorname{Stab}(|\psi\rangle) \subset \mathcal{P}$ can decompose each s_i on Apply H: $X \to Z, Z \to X$ $\{X_1,...,X_N,Z_1,...,Z_N\}$ $\Rightarrow \text{tableau} \left(\begin{array}{cc|c} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{array} \right)$ **Entanglement**: $S(A) = |A| - \log_2 |S_A|$ for bipartition A, B and \mathcal{S}_{A} group of stabilizers acting on A The subgroup \mathcal{S}_A has period two, and therefore $\log_2 |\mathcal{S}_A| \in \mathbb{Z}^{0+}$

zero information $\langle S_{ m tot} \rangle / \mathcal{N}$ 10^{-1} N = 64N = 128policy: select bond i randomly

 \rightarrow dicontinuous phase transition between a volume law and an area law phase \rightarrow divergence of fluctuations at the critical point suggests $p_c \approx 0.38$

complete information VS.

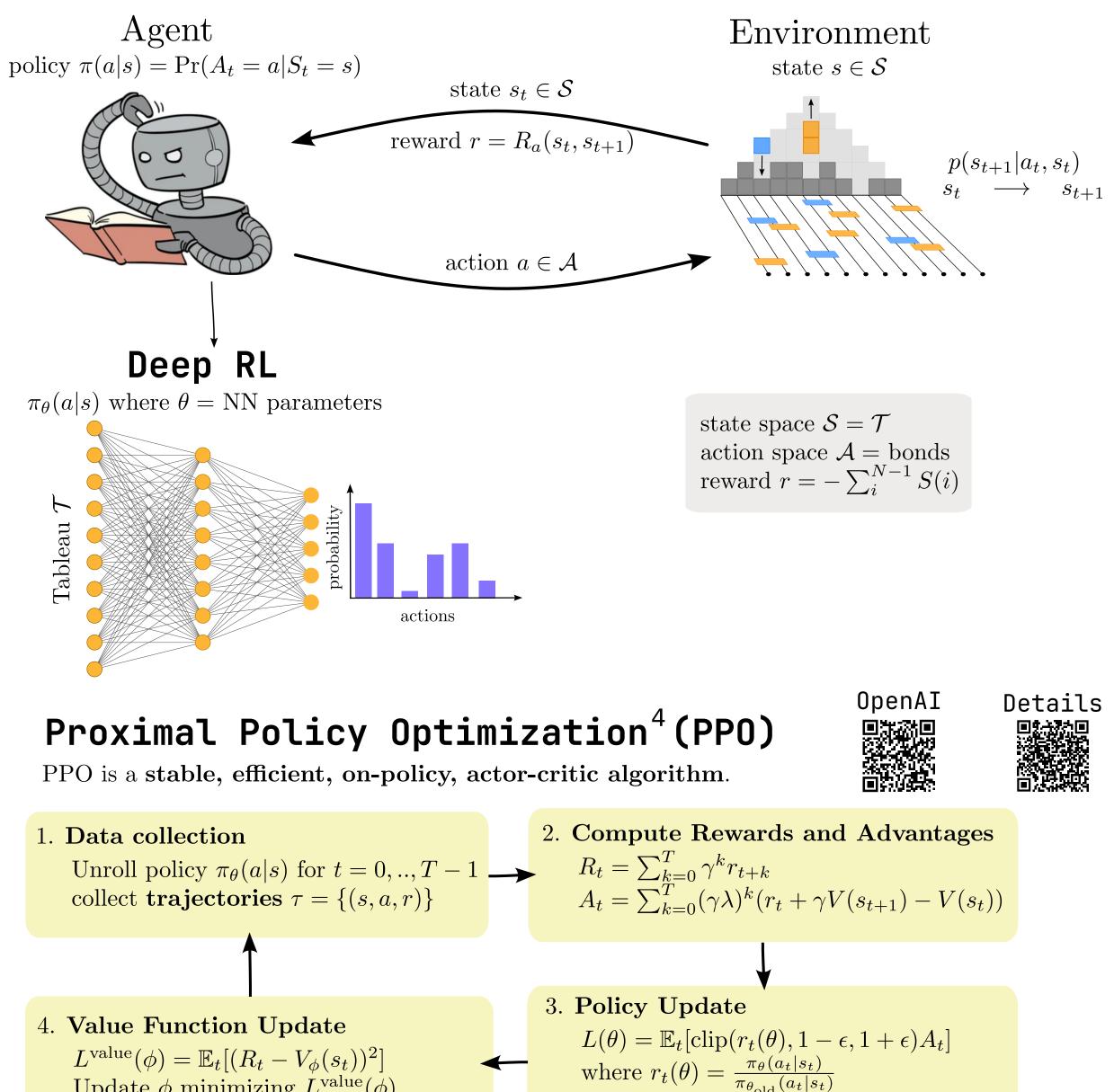


policy: select bond i greedily – bond that maximally disentangles across all i

- \rightarrow continuous phase transition
- \rightarrow divergence of fluctuations at the critical point suggests $p_c \approx 0.15$
- \rightarrow entanglement has a value < 1 for $p \neq 0$

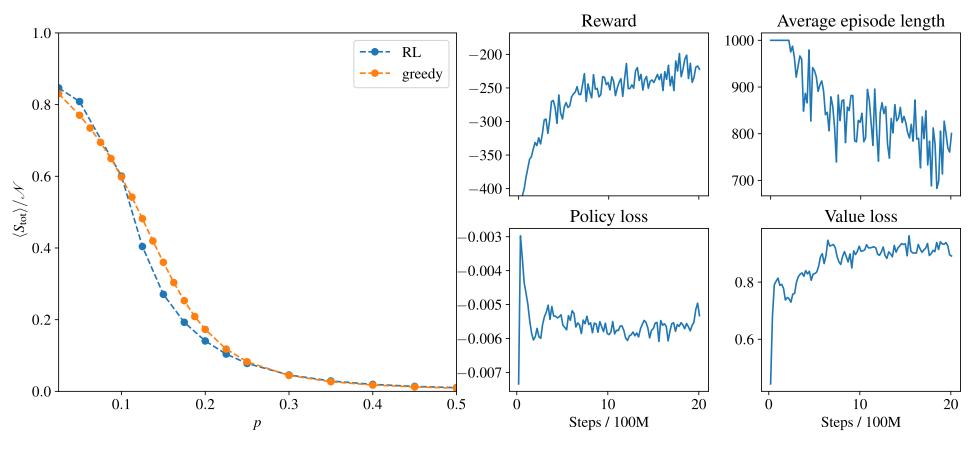
The amount of information available to the Disentanger changes the dynamical properties of the system!

Reinforcement Learning



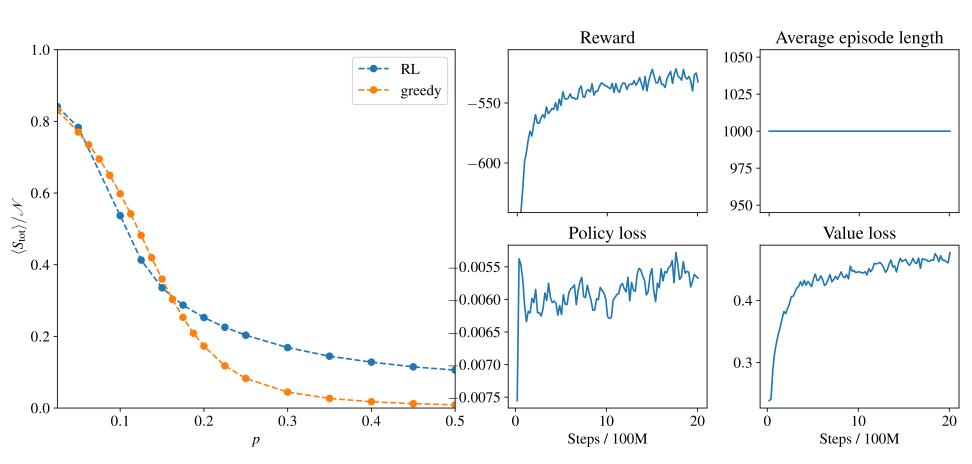
Update θ (i.e. policy) minimizing $L(\theta)$

RL application



Trained at $p = 0.15 \approx p_c$:

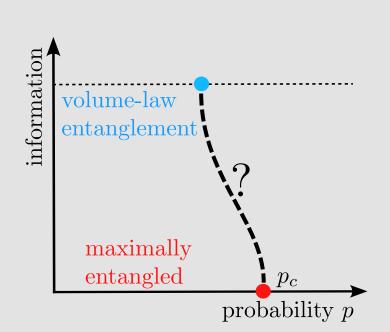
- The RL algorithm converges at a different strategy
- The RL policy outperforms the greedy policy for $p > p^{\text{critical}}$



Trained at $p = 0.1 < p_c$:

- The discrepancy between RL and greedy policy is smaller for $p < p_c$
- Suggests greedy policy is optimal for $p < p_c$

- The greedy strategy is close (but not equal) to the optimal strategy.
- The dynamics of the system is similar in both cases: RL and greedy.
- What is the optimal policy with less information? less information means not all rows of \mathcal{T} are available. There is no straightforward way of generalizing the greedy strategy \rightarrow RL is the right tool for this
- Is complete information necessary to choose optimal local actions? Is there a minimum amount of information s.t. an optimal action can be taken?



We expect two possible scenarios:

- jump from greedy-like to random-like
- continous transformation from greedy-like to random-like

References

Update ϕ minimizing $L^{\text{value}}(\phi)$

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Outlook

- Characterize the dynamical properties at critical point
- Finite size scaling to study behaviour in thermodinamic limit
- Use the RL model to extrapolate between complete and zero information, aiming to study how the phase transition changes

Phase Transitions in Quantum Games

• mpipks

Two agents against each other, each following a *policy*: random, greedy, RL

 \rightarrow who will win, and why? Find out at my poster :)

